

# Hybrid Detection and Mitigation of DNS Protocol MITM attack based on Firefly algorithm with Elliptical Curve Cryptography

Sabitha Banu. A.<sup>1,\*</sup>, Dr. G. Padmavathi<sup>2</sup>

<sup>1</sup> Department of Computer Science, Avinashilingam Institute for Home Science and Higher Education for Women, Coimbatore, India.

<sup>2</sup> Department of Computer Science, Avinashilingam Institute for Home Science and Higher Education for Women, Coimbatore, India. Email: padmavathi.avinashilingam@gmail.com

## Abstract

A Domain Name Server is a critical Internet component. It enables users to surf the web and send emails. DNS is a database used by millions of computers to determine which address best answers a user's query. DNS is an unencrypted protocol that may be exploited in numerous ways. The most popular DNS MITM attack uses DNS poisoning to intercept communications and fake them. DNS servers do not verify the IP addresses they forward traffic to. In DNS attacks, the attacker either targets the domain name servers or attempts to exploit system weaknesses. The Proposed FFOBLA-ECC model detects the DNS Spoofed nodes in a wireless network using the optimized firefly boosted LSTM with the help of TTL and RTR parameters received from the simulation environment and provides authentication between the nodes in order to mitigate it using the Elliptical curve cryptography. The proposed model results are different from the other methods and yield highly accurate results beyond 98% compared with the existing RF, ARF, and KNN methods.

**Keywords:** Domain Name Service(DNS), Man in the Middle attack(MITM), DNS MITM attack, Firefly algorithm, Elliptical Curve Cryptography(ECC).

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## 1. Introduction

The Internet is a haven for all kinds of cyber-attacks. For protection against cyber-attacks, many internet platforms and services designed to make lives more accessible at the same time must be aware and vigilant. Some attacks on web servers and websites have grown simpler due to the Internet's wide accessibility, which causes risks that lead to revenue loss in organizations. One of the most significant risks associated with the cyber-attack is based on the "Domain Name System (DNS)" protocol, which acts as a fundamental internet protocol, DNS Spoofing or DNS MITM (Man in the Middle) attack [1]. The domain name to IP address mapping is made easier with this hierarchical, distributed directory service. DNS Spoofing or DNS MITM attack is where the hackers compromise the

system, inject false DNS records in the DNS servers, and redirect the traffic to some fraudulent websites to steal the credentials or information.

As soon as a domain name is entered into a web browser, the DNS resolver requests to translate the domain name to an IP address. If the domain name isn't discovered in the DNS resolver, then something went wrong. If that's the case, it'll try to find the data on another recursive or iterative server and pass it along to the user so they may follow it. In the DNS server's memory, requests for domain addresses are stored for later use.

DNS is a critical name resolution service that is extensible to clients [2]. However, owing to the enormous increase of internet users, DNS became insecure. Attackers actively seek ways to exploit the system's flaws. DNS cache poisoning, spoofing, or MITM, is a misleading cyber-attack that changes DNS

\* Corresponding author. Email: sabithabanu\_cs@yahoo.com

records via a DNS query that redirects traffic from legitimate servers to false websites. These attacks are difficult to detect and sometimes undetected for a long time, causing serious security issues.

DNS spoofing attacks are classified according to the attacker’s ultimate objective. DNS spoofing may be accomplished in a variety of methods, including the following [3]:

- Compromising a DNS server
- Man in the Middle attack
- DNS Cache Poisoning attack
- Creating False base station network and fabricate the DNS server.

The Global DNS threat report (2019) [4] shows that 82 percent of companies are impacted, with 63 percent suffering outages. The research also showed that financial services companies pay the most per DNS attack compared to other industries.

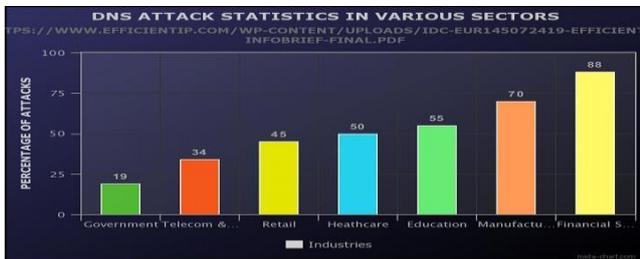


Fig 1 Percentage of DNS Attacks on Various Industries

An example of how DNS Spoofing or DNS MITM works is shown in Fig 2.

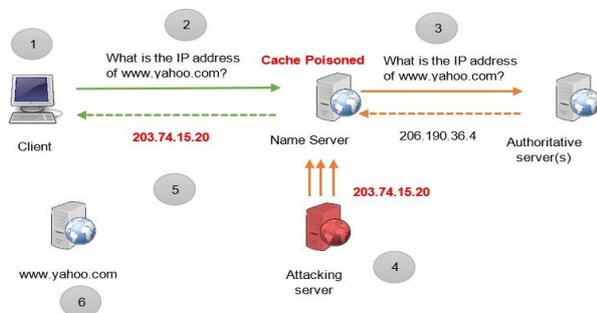


Fig 2 How DNS Spoofing or DNS MITM works.

Different types of DNS attacks have been studied. Among those, DNS MITM or DNS Spoofing are the most predominant and commonly found on the Internet. DNS MITM allows the attackers for data theft, malware infection, stop the security updates and censorship.

The main objective is to predict the detection of the DNS Spoofed nodes in a network using the FFOBLA-ECC algorithm based on the simulation data analysis. The proposed approach creates a network with a server, clients (browsers), and a controller to generate traffic between them. The firefly algorithm optimizes the characteristics based on the generated traffic. The characteristics are given as input into the LSTM model to stabilize it, and then AdaBoost is used to enhance the detection accuracy and provides authentication of the nodes.

The secondary objectives are as follows:

- The Proposed FFOBLA-ECC detection model detected the DNS Spoofed nodes in a wireless environment based on the simulation data without modifying the protocol.
- The predicted accuracy based on the simulation model outperformed the existing Random Forest, Adaptive Random Forest, and KNN by a success rate of 99%.

The rest of the article is structured as follows. Background studies have been discussed in the next part, The FFOBLA-ECC model in part 3 with the detailed protocol and simulation approach. The technical experiment details with the outcomes are discussed in part 4, and Part 5 ends with the final part of the paper.

## 2. Background Study

Some of the necessary online tools that are in practice to handle DNS MITM attacks are

- Ettercap
- Hetty
- Bettercap
- Proxy.py
- Mitmproxy and
- Burp

Based on the literature, different types of Detection mechanisms, Prevention mechanisms, Defense mechanisms, Cryptographic mechanisms, and ML/DL mechanisms to handle the DNS MITM attacks are given in Tables 1,2,3,4 and 5.

The following table 1 summarizes the literature on the various kinds of DNS MITM attack detection methods.

Table 1 DNS MITM attacks Detection mechanisms

Authors	Journal, issue, Year	Detection Mechanism	Limitations
M.Antonakakis et al. [5]	Springer, 2010	Anax(Real Time)	Low rate of False Positives Increased Detection Rate
Yong Wan Ju et al. [6]	IEEE, 2007	Cache Poisoning Detection System(CPDS)( Real Time)	Lower dependency, higher safety, applicability.
Park et al.[7]	ACM Proceedings, 2004, Vol 6	CoDNS	Less secured

L.Poole [8]	ACM Proceedings, 2006, Vol 3	ConfidDNS	Improved security
Sun et al.[9]	Springer, 2009	DepenDNS Protocol	Efficient
Nadhem J. AlFardan, Kenneth G. Paterson [10]	Springer Proceedings, 2010	DepenDNS Protocol	DNS cache poisoning and DDoS attacks went undetected.
L.Yuan et al. [11]	IEEE Proc, 2006, Vol 5	Dox	Simple and modest, improve the coherence of DNS Caches
Roberto Perdisci et al. [12]	Proc, 2008	WSEC DNS (Wild-card Secure DNS)	Overhead in terms of CPU and memory consumption.
Roberto Perdisci et al. [13]	IEEE Proc, 2009	WSEC DNS(Wild-card Secure DNS)	-
B.B.Jaya Singh[14]	CVR Journal of Science and Tech, 2017, Vol 12	Detection Algorithm	-
Sergii Lysenko et al. [15]	Proc, 2020	Isolation Forest algorithm	Blocks the malicious domains based on the DNS traffic
Pratik Satam et al. [16]	Journal of Internet Services and Information Security, 2015, vol no.4	DNS-IDS	
Chirag Sharma [17]	Thesis, 2020	Feed-Forward Neural Network machine learning model	Not applied on the real-time dataset.
Caiyun Huang[18]	IEEE Proc, 2019	Self-Feedback Detection System(SFDS)	High accuracy, effective
Yasuo Musashi et al. [19]	IPJSJ SIG Technical Report, 2011, Vol 53 No 1	Damerau-Lavanshtein Distance based detection model	High accuracy

Yong Wan Ju et al. [20]	IEEE Proc, 2007	CPDS: Cache Poisoning Detection System	Improved trustability
Artem A. Maksutov [21]	IEEE Proc, 2017	DNSwitch	Applied only on LAN
NM Sahri et al. [22]	IEEE Proc, 2016	CAuth	Less computation time, high bandwidth, lightweight method
Hao Wu et al. [23]	IEEE Proc, 2015	Kalman filter	Quick and effective

The following table 2 summarizes the literature on the various kinds of DNS MITM attack Prevention mechanisms.

Table 2 DNS MITM attacks Prevention mechanisms

Authors	Journal, issue, Year	Prevention Mechanisms	Limitations
Nabih Abdelmajid et al. [24]	Proc, 2020	GPS technology for authentication	Drawbacks in Fixed location
Ramzi Bassil et al. [25]	IEEE Proc, 2012	S-DNS	Backward compatible, simple, low communication overhead, low computation overhead, efficient encryption key management scheme.
Jin Cao et al. [26]	Springer, 2017,	Selective Re-Query Case Sensitive Encoding scheme	Increased security, network efficiency.
Haider Salim Hmood et al. [27]	Security in Computer Systems and Networks The Computer Journal, 2015, Vol. 58 No. 4	Adaptive-Cache of DNS (ACDNS)	Accurate enough, reduce latency

Lejun Fan et al. [28]	IEEE Proc, 2011	Security Proxy	Easy to implement, deploy, efficient
Yan Zhao et al. [29]	(JISIS), 2020, vol 10, no 2	DCG client-side protection module	compatible with CDN mechanism, pluggable service
Jayashree Mohan et al. [30]	IEEE Proc, 2015	Bi-Query and S-Key based encoding scheme	Reduces the probability of attack rates
Talha Naqash et al. [31]	IEEE Proc, 2012	Security proxy with SHA-1	Simple to execute
Shimrit Tzur-David et al. [32]	Springer, 2011	Delay Fast Packets	Less memory consumption, scalable, not able to detect DoS attack

The following table 3 summarizes the literature on the various kinds of DNS MITM attack defence mechanisms.

Table 3 DNS MITM attacks Defence mechanisms

Authors	Journal, issue, Year	Defense Mechanisms	Benefits
Zheng Wang et al. [33]	Springer Lec.Notes, 2017, Vol 238	On-Demand Defense (ODD)	Reduces DNSSEC overhead, improves the performance
Tengchao Ma et al. [34]	IEEE Proc,2020	Adaptive dynamic Defense Strategy	Converges to optimization at the local level as opposed to global
Sze Yiu Chau et al. [35]	Springer,2018, Vol 255	CGuard	Easily deployable, compatible
Jonathan Trostle et al. [36]	IEEE Proc,2010	DNS Proxy Server	Minimum performance

The following table 4 summarizes the literature on the various kinds of DNS MITM attack cryptographic methods.

Table 4 DNS MITM attacks Cryptographic mechanisms

Authors	Journal, issue, Year	Cryptographic	Limitations
Mohammed Abdulridha	IEEE Proc, 2016	asymmetric cipher	delay

	issue, Year	Mechanisms	
P.Vixie et al.[37]	IETF, 2000	TSIG	Low-cost to calculate keyed-hash authentication codes
D.E.Eastlake et al.[38]	IETF, 2000	SIG(0)	Expensive asymmetric operations
R.Arends et al. [39]	IETF, 2005	DNSSEC	Do not provide buffer overruns; DoS attacks confidentiality. It can't be deployed in a short period. Increases workload traffic.
D.J.Bernstein [40]	Online, 2009	DNSCurve	No end-to-end security, DNS response validity is not notified.
H.M.Sun et al. [41]	Springer, 2009	S-DNS	Time to Live (TTL) and authoritative answer policies are not adopted. The users accept Outdated SSL certificates and bad certificates.
Naveen Kumar & Kamal Kumar Ranga[42]	IJIRST, 2015, Vol 2, Issue 1	elliptic curve digital signature algorithm	More minor keys with high security, low power consumption.
Manisha Singh & Snigdha [43]	Easy chair, 2020, Preprint no.3793	Message Digest algorithm and RNG(Pseudo Random Number Generator) algorithm, DSA algorithm	
Mohammed Abdulridha	IEEE Proc, 2016	asymmetric cipher	delay

Hussain et al. [44]			
M. H. Jalalzai [45]	IEEE Proc, 2015	cryptographic framework	
Xue Jun Li et al. [46]	IEEE Proc, 2019	character-wise encryption (CWE) method	Reduced encryption time

The following table summarizes the literature of ML/DL techniques to handle DNS MITM attack.

Table 5 Literature of ML/DL Techniques used for DNS MITM attack

Author	Journal, issue, Year	ML /DL Techniques	Limitations
Bai et al.[47]	Springer, 2011	ANN	Accuracy detection-98% (valid and forged packets)
Liguo chen et al.[48]	Elsevier Proceedings, 2018, Vol 134	Random Forest Algorithm	Classification model FPR-0.0% FNR-4.36%
Mykola Kozlenko , Valerii Tkachuk [49]	Proceedings, 2019	RNN	Accuracy detection-70%
Abdallah Moubayed et al. [50]	IEEE Proc, 2020	ensemble-based feature selection and bagging classification model	Lower computational complexity
Cho Do Xuan et al. [51]	IJETER,2020, Vol 8, No 5	Random Forest	Time-consuming on extracting features
Yong Jin et al. [52]	IEEE Proc, 2019	SVM	Not tested in a real-time network.

### 2.1. Observation due to literature

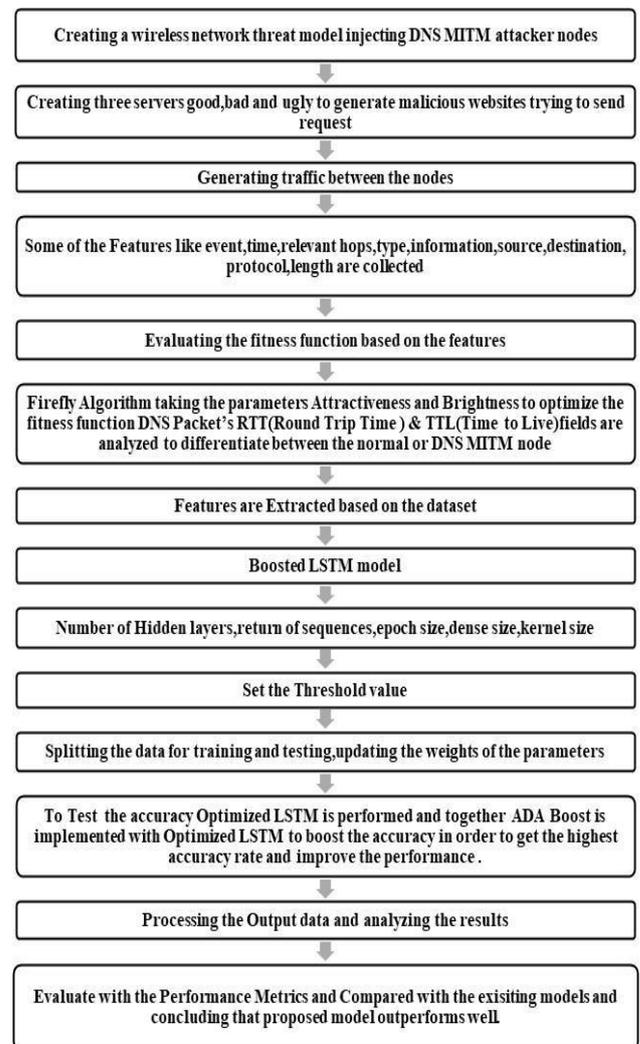
Among all this literature, ML/DL Techniques have given a high accuracy rate where some methods have reduced time in extracting features, some reduced computation time. However, researchers have rare literature on using swarm intelligence algorithms defined in the proposed methodology.

## 3. The Proposed Model

The Proposed FFOBLA-ECC model consists of a server, clients, and a controller which generated the communication between the server and the clients. The model predicts DNS MITM attacker nodes based on the optimized firefly boosted LSTM method, and also it is effective in terms of accurate detection [53].

### 3.1. DNS MITM attacker model

A server, clients (browsers), and a controller are included in the proposed attacker model. The controller generates traffic between the clients and servers. The obtained traffic is captured as an a.csv file, and the firefly algorithm is used to optimize the features. The features are loaded into the LSTM model to stable the model, and the network is then stimulated using AdaBoost to optimize detection accuracy



### 3.2 Protocol Simulation model

An HTTP toolkit with 50 browser clients, three servers, and a single controller is utilized to handle and run the

simulation. In the simulation, the client and server use the TCP/IP model to send messages. Clients and servers generate requests and replies[54]. The scenario includes clients, servers, and two routers linked by a connection. The servers are www.good.com, www.bad.com, and www.ugly.com. The network structure defines the message that may be sent directly between clients and servers or through routers. Some of the browser-based parameters evaluated for the request and reply are given in Table 6.

Table 6. Browser-based Parameters

Parameters
sessionInterval
requestInterval
reqInSession
processingDelay
requestSize
pageSize
numResources
textImageResourceRatio
imageResourceSize
textResourceSize
replyDelay

As part of the attack, the attackers compromised code in the victim’s web browsers, diverting traffic to the attacker’s servers and spoofing the victim’s IP address. As demands increase, so does the server service level. On average, per 200 queries, the browser user receives one malicious link. Even if discovered, the user is routed to the attacker’s site. The firefly optimization method collects browser request and response data over time to identify DNS MITM attacks.

### 3.3. Optimized firefly boosted LSTM Algorithm

The simulation generates datasets in the form of a .csv file. Training and testing datasets are separated. The flashing characteristic of fireflies is the basis for the firefly algorithm [55]. The three rules that apply to flashing features are as follows:

- Fireflies are genderless, so that they may be attracted to each other despite their differences in appearance.
- They are also attractive based on their brilliance; thus, the brighter one attracts the less vivid one.
- The target feature’s search space to be optimized affects the brightness or intensity of a firefly’s light.

Based on the firefly’s brightness, the objective function is optimized on how the brighter firefly is attracted to the other. At the same time, the distance between fireflies enhances attraction. Fireflies move spontaneously if their brightness is the same. A novel

solution created by chance and firefly attraction. Two key issues are at stake: **light intensity variance** and **attractiveness formulation**. A firefly’s brightness  $I$  at a specific location  $x$  defined as  $I(x)$ , and it is proportional to the fitness function  $F I(x)\alpha F(x)$ .

The brightness of the Firefly  $I(r)$  changes with respect to the distance  $r$  may be defined as

$$I = I_0 e^{-\gamma r} \tag{1}$$

Where  $I_0$  is the brightness, and  $\gamma$  is the coefficient of absorption of light. The attractiveness  $\beta$  is equivalent to the brightness of the Firefly and brightness  $\beta$  can be defined as

$$\beta = \beta_0 e^{-\gamma r^2} \tag{2}$$

Where  $\beta_0$  defines the attractiveness at  $r=0$ . Brightness  $I$  and attractiveness  $\beta$  are linked in several ways. Although brightness is a true measure of the firefly’s emitted light, attractiveness is a subjective measure of the light that should be seen in beholders’ eyes and measured by other fireflies. Distance ( $r_{ij}$ ) between the fireflies  $i$  and  $j$  defined using Cartesian distance

$$r_{ij} = \|x_i - x_j\| = \sqrt{\sum_{k=1}^N (x_{ik} - x_{jk})^2} \tag{3}$$

$N$  defines the dimensionality of the problem. The firefly  $i$  migrate towards the brightness of the other firefly  $j$  is defined as

$$x_i = x_i + \beta_0 e^{-\gamma r^2} (x_j - x_i) + \alpha \varepsilon_i \tag{4}$$

Where  $\varepsilon_i$  is a random number that denotes the movement of the fireflies based on the Gaussian distribution. The fireflies’ movement comprises the firefly’s current position, attractiveness towards the other firefly, and random walk a generated randomly from the interval (0, 1).  $\gamma$  denotes the convergence speed. The value is taken from  $\gamma \in (0, \infty)$ .  $\beta_0$  denotes the random walk.

The Firefly method uses the DNS header file’s TTL and roundtrip time to identify DNS MITM attacks. Other features retrieved from the simulation environment include Event, Time, Router, ID, Source got, Source received, Actual received, Actual utilized. Extracted features put into LSTM model to address vanishing gradient. Another kind of recurrent neural network used for time series data prediction LSTM, has several faults, including vanishing gradient issues. It cannot train the model for long-term dependency. LSTM solves these problems by storing data for an extended period utilizing these three gates.

LSTM model uses three gates called input gate, forget gate, and output gate. Each gate learns from the input given through a chain of sequences and chooses whether to retain or reject it to transmit necessary information down the lengthy chain of sequences. The equation of input ( $i_t$ ), output ( $o_t$ ) and forget ( $f_t$ ) states are

$$i_t = \sigma(w_i[h_{t-1}, x_t] + b_i) \tag{5}$$

$$f_t = \sigma(w_f[h_{t-1}, x_t] + b_f) \tag{6}$$

$$o_t = \sigma(w_o[h_{t-1}, x_t] + b_o) \tag{7}$$

$$\tilde{c}_t = \tanh(w_c[h_{t-1}, x_t] + b_c) \tag{8}$$

$$c_t = f_t * c_{t-1} + i_t * \tilde{c}_t \tag{9}$$

$$h_t = o_t * \tanh c_t \tag{10}$$

Where  $\sigma$  denotes the sigmoid function,  $w_x$  denotes the weight of the respective neuron,  $h_{t-1}$  denotes the output received from the earlier iteration at the checksum  $t-1$ ,  $x_t$  denotes entering data of the current checksum,  $b_x$  denotes corresponding gates' prejudices,  $c_t$  denotes cell state memory at the checksum (t),  $\tilde{c}_t$  denotes a possible option for the cell state at the checksum (t). The input gate symbolizes the new information that will be saved in the cell at the moment. The output gate activates the last LSTM block using a sigmoid function, checksum 't' received from the forget state. Forget state represents deleted information. To find a candidate ( $c_t$ ) for the current checksum (t), and from the equations above, the model knows which has to be discarded i.e.,  $f_t * c_{t-1}$  and which has to be taken for the current checksum (t) i.e.,  $i_t * \tilde{c}_t$ . Then the cell state is filtered and passed to the activation function to predict the information that appears in the LSTM block's output in the current checksum (t).  $h_t$  is passed to the current LSTM block with the softmax layer to predict the output  $y_t$ .

When LSTM is trained, the model's performance is continuously unstable. To overcome it, one of the robust ensemble machine learning techniques called AdaBoost, which is good at predicting results, is ensemble with an LSTM block to train iteratively for t iterations.

$$f(x) = \sum_{t=1}^T \alpha_t h_t(x) \tag{11}$$

where

$$\alpha_t = \frac{1}{2} \log \frac{1-s_t}{s_t} \tag{12}$$

$\alpha_t$  denotes the weights of the weak learners in the classifier and combines all the t predictors to boost the model's performance to achieve high accuracy, which results in a stable form. Both AdaBoost and LSTM models are trained separately, and predictions are calculated on average. This combined heterogeneous model yields accurate prediction results than using a single LSTM model.

This hybridized optimized firefly with boosted LSTM gives high accuracy, which outperforms the existing Random Forest, Adaptive Random Forest, KNN models in terms of some performance metrics like Detection Accuracy, Delay, Packet Delivery Ratio, Packet Drop Ratio, Throughput.

**Pseudo Code**

```

1. Identifying Objective function f(x),  $x=(x_1, x_2, \dots, x_d)^T$ 
2. Initialize a population of fireflies  $x_i(i = 1, 2, \dots, n)$ 
3. Define light absorption coefficient gamma
4. WHILE count < MaximumGenerations
5. FOR  $i = 1 : n$  (all n fireflies)
6. FOR  $j = 1 : i$ 
7. Light intensity  $I_i$  at  $x_i$  is determined by  $f(x_i)$ 
8. IF  $I_i > I_j$ 
9. Move firefly i towards j in all d dimensions
10. ELSE
    
```

```

11. Move firefly i randomly
12. END IF
13. Attractiveness changes with distance r via  $\exp[-\gamma r^2]$ 
14. Determine new solutions and revise light intensity
15. END FOR j
16. END FOR i
17. Rank the fireflies according to light intensity and find the current best
18. END WHILE
19. Define Features_Cols(Event,Time,Router,ID,Source got,Source Received,Actual Received,Actual Used)
20. Assign  $x=feature\_Cols, y=label$ 
21. Define  $X\_train, X\_test, y\_train, y\_test$ 
22. Define  $num\_words = 1150892$ 
23. Define LSTM model = Sequential()
24. Set LSTM model(Embedding(num_words, 6, input_length=8))
25. Set LSTM model.add(Dropout(0.1))
26. Set LSTM model(25, dropout=0.1, recurrent_dropout=0.1))
26 Set LSTM model(Dense(75, activation='relu'))
27. Set LSTM model(Dropout(0.1))
28. Set LSTM model.add(Dense(1, activation='sigmoid'))
29. Set LSTM model.compile(loss='binary_crossentropy', optimizer='adam', metrics=['accuracy'])
30. Set LSTM model.fit(X_train, y_train, batch_size=54, epochs=10)
31. Print Detection Accuracy.
32. Stop
    
```

### 4. Experiment Results And Discussions

The Technical specifications of the experiment conducted in a simulated environment are given in table 7. Six months of data logs were taken for the experiment. The Proposed hybrid FFOBLA-ECC model gives a novel solution to detect and mitigate the DNS MITM attack nodes using the Optimized firefly boosted LSTM and ECC algorithm. It produced very accurate outcomes when compared with the existing RF, ARF, and KNN models. The advantages of this proposed hybrid FFOBLA-ECC model include its ease of convergence into complicated problems, its simplicity and versatility, and its accurate, high-performance outcomes.

Table 7. Simulation Parameters

Parameters	Values
Number of Access requests	50
Initial energy level	100(j)
Number of Objects	50
Energy consumption for reception and transmission	.395(w),660(w)
Zone of simulation	1500m×1500m
Transmission Range	100kbps
Number of Services	50
Simulation time	3600s

Simulation is done, and the data are collected, and the packet details are

- Event
- Time
- Relevant Hops
- Type
- Information
- Source

- Destination
- Protocol
- Length

Performance measures are used to compare the Existing and Proposed approaches. Additionally, they are

- Detection Accuracy
- Delay
- Packet Delivery Ratio
- Packet Drop Ratio
- Throughput

### 4.1. Detection Accuracy Ratio

The proportion of correctly predicted DNS MITM nodes is about the total number of predictions produced. It can be shown from the detection accuracy ratio that the hybrid FFOBLA-ECC model outperforms the current methods.

$$DR = \frac{TP + TN}{TP + TN + FP + FN}$$

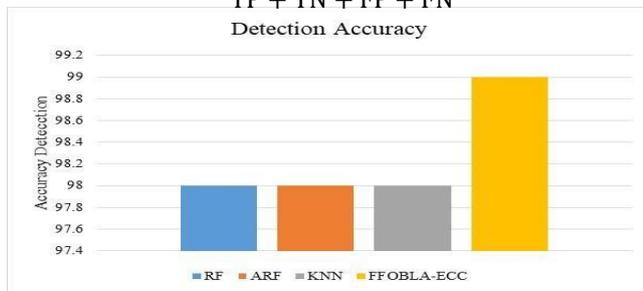


Fig 3 Accuracy Detection Ratio for DNS MITM Nodes

The given graph demonstrates that the proposed hybrid model accurately computed the DNS MITM attack nodes when applied to the simulation dataset. Existing techniques such as Random Forest, Adaptive Random Forest, and KNN have shown detection accuracy of 98%, 98%, and 98%, respectively. The Proposed hybrid model, compared to other methods, the optimized firefly boosted LSTM algorithm outperforms them all and obtained a detection accuracy of 99%.

### 4.2. Average Delay

The time it takes for data packets to arrive at their destination. Furthermore, it takes time to find the routes and queue packets for transmission. It only counts packets that are delivered to their destination. When the average Delay outcomes are low, it indicates better performance. Average delay is shown in Fig 4.

$$\text{Average Delay} = \frac{\text{Total no of arriving time} - \text{Total no of sending time}}{\text{Total no of Packets}}$$

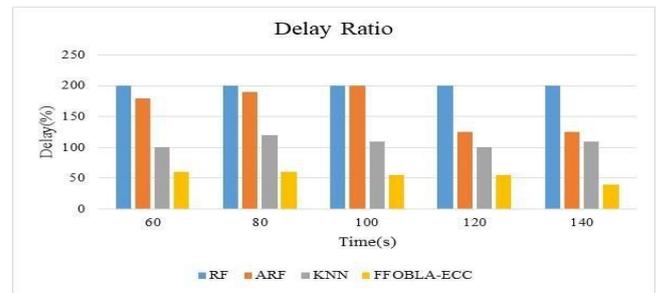


Fig 4 Average Delay

### 4.3. Packet Delivery Ratio

The source node's and the destination node's combined packet counts. The mathematical equation used to calculate PDR is shown below.

$$\text{Packet Delivery Ratio} = \left( \frac{\text{Packets Received}}{\text{Packets Sent}} \right) \times 100\%$$

Packet Delivery Ratio Fig 5 is given below.

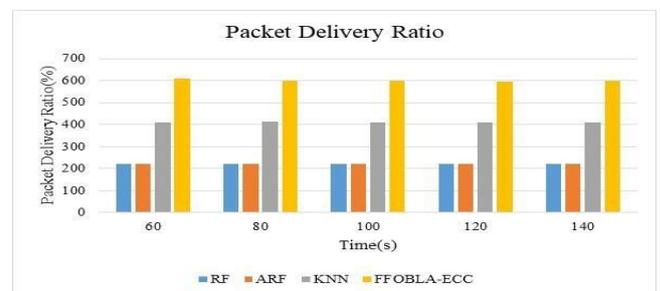


Fig. 5 Packet Delivery ratio

### 4.4. Packet Drop Ratio

Packet loss is a percentage of overall packet transmission time. It is calculated by Packet drop ratio(%) = sent packets – received packets. Packet drop ratio Fig 6 is given below.

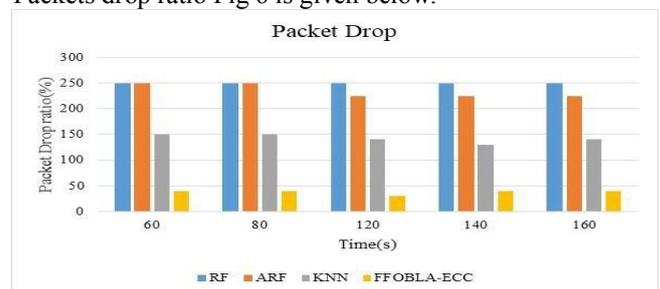


Fig 6 Packet drop Ratio

### 4.5. Throughput

Throughput is the number of packets successfully acquired in a given period, measured in bits per second (bps). Throughput Fig 7 given below

$$\text{thruput} = \sum \frac{(\text{Traffic received} - \text{Traffic sent})}{\text{Total data packets received}}$$

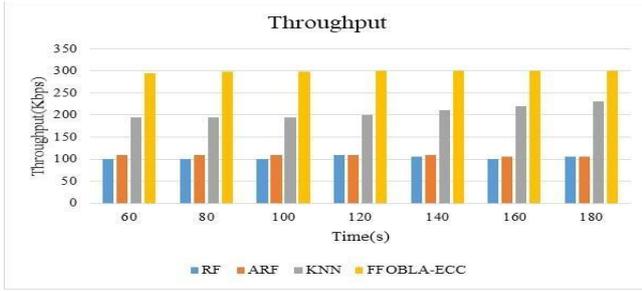


Fig 7 Throughput

Table 7.7 compares the proposed FFOBLA-ECC model to existing methods in terms of performance measures. The proposed hybrid FFOBLA-ECC Method performed optimally based on time (s) basis as well as the total number of nodes, resulting in a small percentage of delay. Packet Delivery Ratio was the most outstanding level reached 600 percent compared to previous techniques that attained just 200 percent, resulting in packet delivery success. The packet drop percentage is low, ranging between 30% and 40%. The current techniques RF, ARF, and KNN, had the most significant drop ratios ranging from 150 to 250 percent, but the maximum throughput ratio of 300 percent remained constant. All of the time and current techniques range from 100 percent to 225 percent, resulting in the best packet transmission.

Table 7 Comparison of Existing Model and Proposed Model

Metrics used for comparison	Existin g: RF (bps)	Existin g: ARF (bps)	Existin g: KNN (bps)	Propose d: FFOBLA -ECC (bps)
Detection Accuracy	98	98	98	99
Average Delay	200	200	100	60
Packet Delivery Ratio	220	220	410	610
Packet Drop Ratio	250	250	150	40
Throughput	100	110	195	295

### 5. Conclusion

The DNS MITM attack is where the attacker holds the DNS records and tries to change them so that it can redirect the traffic to fake websites to steal the credentials of the victims or access some sensitive

information or malicious websites try to install worms or virus software in their personal computer for long term access to the data stored in the computer.

DNS MITM attack can be troublesome for both the website owners and the users. It can cause some security issues and be undetected for an extended period. There are various prevention techniques applied to avoid DNS MITM attacks using some encryption techniques, DNSSEC. Various literature was given to detect, prevent, and defend against the DNS MITM attack. The Proposed hybrid Model uses a threat model of creating a DNS MITM attack using a simulation tool and injecting the attack, capturing the traffic before and after the attack.

Using Optimized Fireflies boosted LSTM algorithm provides a novel solution in detecting and predicting the attacker nodes' accuracy by using the captured file obtained from the simulation. Based on the TTL and Roundtrip time of the DNS record, some features are extracted and fed into the LSTM model for the training. It uses AdaBoost with the LSTM model to stabilize the network and improve its performance to achieve high accuracy results. Boosted LSTM performs well than the single LSTM model.

The Proposed FFOBLA-ECC Model outperforms the existing model and yields a high accuracy rate with performance.

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